# Structure of Incoherent Operations 

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## Structure of the Resource Theory of Quantum Coherence

Alexander Streltsov, Swapan Rana, Paul Boes, and Jens Eisert
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 mbed-state conversion problem for a single qubit

## Outline

- A glance at QRT of coherence
- Mathematical framework for superposition principle
- Many models for free operations
- Structure of Incoherent operations
- General quantum operations
- Number of Kraus Operators and its importance
- Qubit incoherent channel
- Improved bound on \# Kraus operators for IO
- Exact number for SIO
- Achievable region and collapse of hierarchies
- Application to quantum thermodyanmics


## Coherence at a glance

Many model of Coherence theory [Streltsov et al., Rev. Mod. Phys. (2017)]
$\left.\begin{array}{cccccc}\hline & 1 & 2 & 3 & 4 & 5 \\ \hline \text { MIO } & \text { (Åberg, 2006) } & \text { yes yes yes yes yes } \\ \text { IO } & \begin{array}{c}\text { (Baumgratz et al., 2014; Winter } \\ \text { and Yang, 2016) }\end{array} \\ \text { yes yes yes yes yes }\end{array}\right]$

Table II List of alternative frameworks of coherence with respect to our criteria 1-6 provided in the text.

## Resource Theory of Quantum Coherence

## IO theory of coherence [Baumgratz et al., PRL (2014)]

2Free (Incoherent) states: Diagonal states $\delta=\sum \delta_{i}|i\rangle\langle i|$, for a preferred/chosen o.n.b. $\{|i\rangle\}$. This is not a shortcoming!
2. Free (Incoherent) operations: $\Lambda$ is incoherent iff there is a Kraus decomposition $\Lambda=\left\{K_{n}\right\}$ such that $K_{n} \delta K_{n}^{\dagger}$ is diagonal for all $\delta, n$.

- Maximally coherent state: $\left|\Phi_{d}\right\rangle=\frac{1}{\sqrt{d}} \sum|i\rangle$.
- Any $\rho \in \mathscr{B}\left(\mathscr{H}^{d}\right)$ can be created from $\left|\Phi_{d}\right\rangle$ :

$$
\left|\Phi_{d}\right\rangle \xrightarrow[\text { with certainty }]{\text { only } \Lambda \in \mathscr{I}} \rho
$$

- $\left|\Phi_{d}\right\rangle$ allows to implement arbitrary unitary $U \in S U(d)$.
- Existence of $\left|\Phi_{d}\right\rangle$ allows all kind of concepts related to manipulation of resource e.g., formation, cost, distillation etc.

Quantum operations


## Quantum operations



Quantum operations


## Quantum Operations on a single system

- Are described by Maps: $\rho^{\prime}=\varepsilon(\rho)$
- Two simplest examples: Unitary Evolution $\rho \mapsto U \rho U^{\dagger}$, Measurement $\rho \mapsto \rho_{m}:=K_{m} \rho K_{m}^{\dagger} / \operatorname{Tr}\left[K_{m} \rho K_{m}^{\dagger}\right]$.
- Quantum operations $=$ Quantum channels $=$ CPTP maps

$$
\begin{aligned}
\rho^{\prime} & =\varepsilon(\rho) \\
& =\operatorname{Tr}_{E}\left[U\left(\rho \otimes|0\rangle_{E}\langle 0|\right) U^{\dagger}\right] \\
& =\sum_{m}\langle m| U\left(\rho \otimes|0\rangle_{E}\langle 0|\right) U^{\dagger}|m\rangle \\
& =\sum_{m} K_{m} \rho K_{m}^{\dagger}, \quad K_{m}=\langle m| U|0\rangle \in \mathscr{B}\left(\mathscr{H}^{S}\right) .
\end{aligned}
$$

- $K_{m}$ 's are known as Kraus operators, completely describe the action of the map/channel.
- The $\left\{K_{m}\right\}$ s is in general not unique: Two sets $\left\{K_{m}\right\}$ and $\left\{L_{n}\right\}$ generate the same channel iff

$$
K_{m}=\sum_{n} U_{m n} L_{n}, \quad \forall m, n
$$

Here $U_{m n}$ is a unitary matrix of order $\max \{m, n\}$. This follows essentially from the same result for ensemble:

$$
\left\{p_{i},\left|\psi_{i}\right\rangle\right\}=\left\{q_{j},\left|\phi_{j}\right\rangle\right\} \text { iff } \sqrt{p_{i}}\left|\psi_{i}\right\rangle=\sum_{j} U_{i j} \sqrt{q_{j}}\left|\phi_{j}\right\rangle
$$

- This implies: If $\rho \in \mathscr{B}\left(\mathscr{H}^{d}\right), \#\left(K_{m}\right) \leq d^{2}$.
- Thus a qubit channel can be described by at most 4 Kraus operators.
- However, we don't know which $U$ will give us incoherent $\Lambda$.

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- Important for simulating $\Lambda$.
- On qubit level, allows to visualize all possible $\Lambda[\rho]$.
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Question: How many Kraus operators are needed to write $\Lambda_{d} \in \mathscr{I}$ ?

- Important for simulating $\Lambda$.
- On qubit level, allows to visualize all possible $\Lambda[\rho]$.
- We have only partial answers :
- General upper bounds for IO and SIO
- Exact results for qubits only


## Upper bound from Choi-Jamiołkowski+Caratheodory

## Upper bound on \#Kruas operators for IO channel

Any incoherent operation acting on a Hilbert (state) space of dimension $d$ admits a decomposition with at most $d^{4}+1$ incoherent Kraus operators.

- Choi-Jamiołkowski isomorphism between a quantum operation $\Lambda$ and the corresponding Choi state :

$$
\rho_{\Lambda}=(\Lambda \otimes \mathbb{1})\left(\Phi_{d}^{+}\right), \quad \Phi_{d}^{+}=d^{-1} \sum_{i, j=0}^{d-1}|i, i\rangle\langle j, j|, \operatorname{dim}\left(\Phi_{d}^{+}\right)=d^{2} .
$$

The rank of the Choi state is the Kraus rank, which is the smallest number of (not necessarily incoherent) Kraus operators.

- Consider the operator

$$
M=(K \otimes \mathbb{1}) \Phi_{d}^{+}\left(K^{\dagger} \otimes \mathbb{1}\right) \text { with any incoherent } K
$$

For any incoherent operation $\Lambda$, the corresponding Choi state $\rho_{\Lambda}$ belongs to the convex hull of the operators $M$. Applying Caratheodory on $M$ gives the upper bound.

## Qubit IO channel

## \#Kraus operators $\leq 5$ for IO

Any qubit IO channel $\Lambda$ admits a decomposition with at most 5 incoherent Kraus operators. A canonical choice of the operators is given by the set

$$
\left\{\left(\begin{array}{cc}
a_{1} & b_{1} \\
0 & 0
\end{array}\right),\left(\begin{array}{cc}
0 & 0 \\
a_{2} & b_{2}
\end{array}\right),\left(\begin{array}{cc}
a_{3} & 0 \\
0 & b_{3}
\end{array}\right),\left(\begin{array}{cc}
0 & b_{4} \\
a_{4} & 0
\end{array}\right),\left(\begin{array}{cc}
a_{5} & 0 \\
0 & 0
\end{array}\right)\right\}
$$

where $a_{i}$ can be chosen $\geq 0$, while $b_{i} \in \mathbb{C}$. Moreover, it holds that $\sum_{i=1}^{5} a_{i}^{2}=\sum_{j=1}^{4}\left|b_{j}\right|^{2}=1$ and $a_{1} b_{1}+a_{2} b_{2}=0$.

- The incoherent condition implies that the Kraus operators can have at most non-zero element in a column.
- Group them into four categories:

$$
\begin{array}{rlrl}
K^{I} & =\left\{\left(\begin{array}{cc}
* & * \\
0 & 0
\end{array}\right)\right\}, & K^{I I}=\left\{\left(\begin{array}{cc}
* & 0 \\
0 & *
\end{array}\right)\right\}, \\
K^{I I I}=\left\{\left(\begin{array}{ll}
0 & 0 \\
* & *
\end{array}\right)\right\}, & K^{I V}=\left\{\left(\begin{array}{cc}
0 & * \\
* & 0
\end{array}\right)\right\} .
\end{array}
$$

- The unitary equivalence $L_{i}=\sum_{j} U_{i, j} K_{j}$ reduces them to eight

$$
\begin{aligned}
K^{I}=\left\{\left(\begin{array}{ll}
* & 0 \\
0 & 0
\end{array}\right),\left(\begin{array}{ll}
* & * \\
0 & 0
\end{array}\right)\right\}, & K^{I I}=\left\{\left(\begin{array}{ll}
* & 0 \\
0 & 0
\end{array}\right),\left(\begin{array}{ll}
* & 0 \\
0 & *
\end{array}\right)\right\}, \\
K^{I I I}=\left\{\left(\begin{array}{ll}
0 & 0 \\
* & 0
\end{array}\right),\left(\begin{array}{ll}
0 & 0 \\
* & *
\end{array}\right)\right\}, & K^{I V}=\left\{\left(\begin{array}{ll}
0 & 0 \\
* & 0
\end{array}\right),\left(\begin{array}{cc}
0 & * \\
* & 0
\end{array}\right)\right\} .
\end{aligned}
$$

- Gathering altogether leads to six Kraus operators:

$$
\left\{\left(\begin{array}{ll}
* & * \\
0 & 0
\end{array}\right),\left(\begin{array}{ll}
0 & 0 \\
* & *
\end{array}\right),\left(\begin{array}{ll}
* & 0 \\
0 & *
\end{array}\right),\left(\begin{array}{ll}
0 & * \\
* & 0
\end{array}\right),\left(\begin{array}{ll}
* & 0 \\
0 & 0
\end{array}\right),\left(\begin{array}{ll}
0 & 0 \\
* & 0
\end{array}\right)\right\} .
$$

- Among these consider the following 3 operators

$$
K_{1}=\left(\begin{array}{cc}
0 & 0 \\
a_{1} & b_{1}
\end{array}\right), K_{2}=\left(\begin{array}{cc}
a_{2} & 0 \\
0 & b_{2}
\end{array}\right), K_{3}=\left(\begin{array}{cc}
0 & 0 \\
a_{3} & 0
\end{array}\right) .
$$

The unitary

$$
U=\left(\begin{array}{ccc}
l a_{1}^{*} & 0 & l a_{3}^{*} \\
m b_{1}^{*}\left|a_{3}\right|^{2} & m\left(\left|a_{1}\right|^{2}+\left|a_{3}\right|^{2}\right) b_{2}^{*} & -m a_{3}^{*} b_{1}^{*} a_{1} \\
n a_{3} b_{2} & -n a_{3} b_{1} & -n a_{1} b_{2}
\end{array}\right)
$$

transforms those to

$$
L_{1}=\left(\begin{array}{cc}
0 & 0 \\
* & *
\end{array}\right), L_{2}=\left(\begin{array}{cc}
* & 0 \\
0 & *
\end{array}\right), L_{3}=\left(\begin{array}{cc}
* & 0 \\
0 & 0
\end{array}\right)
$$

- Thus, altogether those reduces to the following five

$$
\left\{\left(\begin{array}{ll}
* & * \\
0 & 0
\end{array}\right),\left(\begin{array}{ll}
0 & 0 \\
* & *
\end{array}\right),\left(\begin{array}{cc}
* & 0 \\
0 & *
\end{array}\right),\left(\begin{array}{cc}
0 & * \\
* & 0
\end{array}\right),\left(\begin{array}{cc}
* & 0 \\
0 & 0
\end{array}\right)\right\} .
$$

The canonical parameterization follows from the completeness relation $\sum K_{i}^{\dagger} \cdot K_{i}=1$.
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$$
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0 & 0 \\
* & *
\end{array}\right), L_{2}=\left(\begin{array}{cc}
* & 0 \\
0 & *
\end{array}\right), L_{3}=\left(\begin{array}{cc}
* & 0 \\
0 & 0
\end{array}\right)
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* & * \\
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\end{array}\right),\left(\begin{array}{cc}
0 & 0 \\
* & *
\end{array}\right),\left(\begin{array}{cc}
* & 0 \\
0 & *
\end{array}\right),\left(\begin{array}{cc}
0 & * \\
* & 0
\end{array}\right),\left(\begin{array}{cc}
* & 0 \\
0 & 0
\end{array}\right)\right\} .
$$

The canonical parameterization follows from the completeness relation $\sum K_{i}^{\dagger} \cdot K_{i}=1$.

Update: A generic qubit IO $\Lambda$ can be decomposed into four Kraus operators. We are trying to prove that four is indeed the optimal number.

## Exact number for qubit SIO is four

- A canonical form for any qubit SIO is given by

$$
\left\{\left(\begin{array}{cc}
a_{1} & 0 \\
0 & b_{1}
\end{array}\right),\left(\begin{array}{cc}
0 & b_{2} \\
a_{2} & 0
\end{array}\right),\left(\begin{array}{cc}
a_{3} & 0 \\
0 & 0
\end{array}\right),\left(\begin{array}{cc}
0 & 0 \\
a_{4} & 0
\end{array}\right)\right\},
$$

where $a_{i} \geq 0$ and $\sum_{i=1}^{4} a_{i}^{2}=\sum_{j=1}^{2}\left|b_{j}\right|^{2}=1$.

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0 & b_{2} \\
a_{2} & 0
\end{array}\right),\left(\begin{array}{cc}
a_{3} & 0 \\
0 & 0
\end{array}\right),\left(\begin{array}{cc}
0 & 0 \\
a_{4} & 0
\end{array}\right)\right\}
$$

where $a_{i} \geq 0$ and $\sum_{i=1}^{4} a_{i}^{2}=\sum_{j=1}^{2}\left|b_{j}\right|^{2}=1$.

## Bound for higher ( $d$-) dimensional channels

- IO: $\# \leq d\left(d^{d}-1\right) /(d-1)$. Better than $d^{4}+1$ only for $d \leq 3$.
- SIO: $\# \leq \sum_{k=1}^{d} d!/(k-1)$ !. Better than $d^{4}+1$ only for $d \leq 5$.
- (S)IO: \# $\geq d^{2}$ as the set of standard matrix units are linearly independent and forms an (S)IO.


## Application: Achievable region for qubit

## $\mathrm{SIO}=\mathrm{IO}=\mathrm{MIO}$



Figure: Achievable region for single-qubit SIO, IO, and MIO. Colored areas show the projection of the achievable region in the $x-z$ plane for initial Bloch vectors $(0.5,0,0.5)^{T}$ [blue], $(-0.8,0,-0.6)^{T}$ [green], and $(1,0,0)^{T}$ [red]. Note that the last two states are pure. The magenta line corresponds to the achievable region of an incoherent state with Bloch vector $(0,0,0.65)^{T}$.

## Quantum Thermodynamics: Gibbs-preserving SIO

- Any $p|0\rangle\langle 0|+(1-p)|1\rangle\langle 1| \in \mathscr{I}$ can be interpreted as a Gibbs state $\tau=e^{-\beta H} / \operatorname{Tr}\left[e^{-\beta H}\right]$, for a suitable inverse temperature $\beta=\frac{1}{k T}$ and Hamiltonian $H$ which is diagonal,

$$
p=\frac{e^{-\beta E_{0}}}{e^{-\beta E_{0}}+e^{-\beta E_{1}}} .
$$

- Thermal operations are Gibbs-preserving: $\Lambda[\tau]=\tau$, but can create coherence.


Figure: Achievable region [blue area] of single-qubit SIO which preserve the state $t=(0,0,1)$ [left figure] and $t=(0,0,-1)$ [right figure]. The initial state has the Bloch vector $r=(0.5,0,0.5)^{T}$ [blue dot], and the corresponding Bloch vector $\boldsymbol{t}$ is shown as a green dot.

## Conclusion and Outlook

- On qubit level, any SIO or a generic IO can be decomposed into four Kraus operators. This significantly reduces the number of parameters to simulate those channels, as well as to find the exact achievable regions for a given input states.
- The bound on number of Kraus operators derived from combinatorial arguments gives better result in small dimension only. There must be some further unitary reductions.
- We conjecture that every qubit IO could be decomposed into four Kraus operators.
- The restrictions on Kraus operators to define free operations is too strong which has lead to so many RTQCs. There is probably a deeper question involved: if the Krus operators are restricted to have a (sparse) pattern, then how to efficiently bound their number? How to physically implement those operations?


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EXCELENCIA SEVERO OCHOA

